Benha University
$1^{\text {st }}$ Term (January 2019) Final Exam
Class: 4th Year Students (Computer Science Major)
Subject: Compiler Theory
Course Code: CSW 456

Faculty of Computers \& Informatics
Date: 10/1/2019
Time: 3 Hours
Total Marks: 75 Marks
Examiner(s): Dr. Ahmed Hassan

## Answer the following questions [ $\mathbf{6}$ questions in $\mathbf{2}$ pages):

## Question No. 1

## (a)Define the following:

(1) Lexemes : is a sequence of characters in the source program that matches the pattern for a token and is identified by the lexical analyzer as an instance of that token.
(2) Lexical Analyzer : read the input characters of the source program. Group them into lexemes. Produce as output a sequence of tokens for each lexeme in the source program.
(3) Parser : It takes the token produced by lexical analysis as input and generates a parse tree (or syntax tree). In this phase, token arrangements are checked against the source code grammar.
(b)Draw the phases of a compiler.

(a) Find a regular expression for the language of all strings over $\{\mathbf{a}, \mathrm{b}\}$ with odd number of "a" and ending with abb.
$b^{*}(a b * a) * b^{*} \mathbf{a b b} \quad$ OR $(b+a b * a) * a b b$ OR any one equivalent.
(b) convert the regular expression " $\mathbf{a b}+(\mathbf{a}+\mathbf{b})^{*}$ " to NFA.


Question No. 3

$$
\begin{aligned}
& \mathbf{S} \rightarrow \mathbf{S}+\mathbf{S}|\mathbf{S}-\mathbf{S}| \mathbf{T} \\
& \mathbf{T} \rightarrow \mathbf{S} * \mathbf{T}|\mathbf{S} / \mathbf{T}| \mathbf{a}
\end{aligned}
$$

(a) Prove that the grammar is ambiguous.

(b) Remove the left factor then the left recursion from the grammar.

| $\mathbf{S} \rightarrow \mathbf{S A} \mid \mathbf{T}$ | $\mathrm{S} \rightarrow \mathrm{TE}$ | $\mathrm{S} \rightarrow \mathrm{TE}$ |
| :---: | :---: | :---: |
| $\mathbf{A} \rightarrow+\mathbf{S} \mid-\mathbf{S}$ | $\mathbf{E} \rightarrow \mathbf{A E} \mid \boldsymbol{E}$ | $\mathbf{E} \rightarrow \mathbf{A E} \mid \boldsymbol{E}$ |
| $\mathbf{T} \rightarrow \mathbf{S B} \mid \mathbf{a}$ | T $\rightarrow$ TEB \| a | $\mathrm{T} \rightarrow \mathbf{a F}$ |
| $\mathrm{B} \rightarrow * \mathrm{~T} \mid / \mathrm{T}$ | $\mathbf{A} \rightarrow+\mathbf{S} \mid-\mathbf{S}$ | $\mathbf{F} \rightarrow \mathbf{E B F}\|\boldsymbol{\varepsilon} \Rightarrow \mathrm{F} \rightarrow \mathrm{AEBF}\| \mathcal{E}$ |
|  | $\mathrm{B} \rightarrow * \mathrm{~T} \mid / \mathrm{T}$ | $\mathbf{A} \rightarrow+\mathbf{S} \mid-\mathrm{S}$ |
|  |  | $\mathbf{B} \rightarrow * \mathbf{T} \mid$ T |



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## A language contains three types of tokens as the following:

## 1-The first token type is the keword = \{ if \}, (higher priority)

2-The second token type represents the identifiers which are any non-empty string over $\{\mathbf{a}, \mathbf{i}, \mathbf{f}\}$,
$3-$ The third token type are the unary integers $=\{1,11,111,1111, \ldots\}$.
(lower priority)

## For the given language do:

(a) Write a pattern (regular expression) to define the lexemes of each token.

$$
\text { if } \quad(a|i| f)(a|i| f)^{*} \quad 11^{*}
$$

(b)Draw an NFA scanner for your patterns from step (a).

(c) Transform the NFA scanner from step (b) into DFA scanner.


| state | $\mathbf{i}$ | $\mathbf{f}$ | $\mathbf{a}$ | 1 |
| :---: | :---: | :---: | :---: | :---: |
| 0123 | 45 | 5 | 5 | 6 |
| 45 ID | 5 | 57 | 5 | $\varnothing$ |
| 5 ID | 5 | 5 | 5 | $\varnothing$ |
| 57 IF | 5 | 5 | 5 | $\varnothing$ |
| 6 INT | $\varnothing$ | $\varnothing$ | $\varnothing$ | 6 |

Use the scanner from step (c) to define the tokens types and lexemes in the following input stream " ifaa111if11biif11"

$$
\begin{gathered}
\text { ifaa } \rightarrow \text { ID } \\
111 \rightarrow \text { INT } \\
\text { if } \rightarrow \text { IF } \\
11 \rightarrow \text { INT } \\
\mathbf{b} \rightarrow \text { Error } \\
\text { iif } \rightarrow \text { ID } \\
11 \rightarrow \text { INT }
\end{gathered}
$$

Question No. 5
For the following grammar:

$$
\begin{gathered}
\mathbf{S} \rightarrow \mathbf{a S b}|\mathbf{a A}| \mathbf{a} \\
\mathbf{A} \rightarrow \mathbf{c}
\end{gathered}
$$

(a) Construct the LR(1) parser table.

(b) Use the parser table to check the string: aacb

| Stack | Input | Action |
| :--- | ---: | :---: |
| \$0 | aacb\$ | S3 |
| \$0a3 | acb\$ | S3 |
| \$0a3a3 | $\mathrm{cb} \$$ | S5 |
| \$0a3a3c5 | b\$ | r4 and goto 4 |
| \$0a3a3A4 | b\$ | r2 and goto 2 |
| \$0a3S2 | b\$ | S6 |
| \$0a3S2b6 | \$ | r1 and goto 1 |
| \$0S1 | \$ | accepted |



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(c) Is it LR(0) grammar? Why? No, state $\mathbf{3}$ has both shift and reduce.

Question No. 6
Use the following semantic rules to generate the intermediate code for
3*4+5
And what is the result?


| INDEX | VALUE |
| :--- | :--- |
| $\mathbf{0}$ | Int 3 |
| $\mathbf{1}$ | Int 4 |
| $\mathbf{2}$ | Int 5 |
| $\mathbf{3}$ | $\mathbf{1}+\mathbf{2}$ |
| $\mathbf{4}$ | $\mathbf{0} * \mathbf{3}$ |

$\mathrm{T} 0=3$
$\mathrm{T} 1=4$
T2 $=5$
T3 $=\mathbf{T} 1+\mathrm{T} 2$
$\mathrm{T} 4=\mathrm{T} 0 * \mathrm{~T} 3$
Result $=3 *(4+5)=3 * 9=27$

